Quantifying Information Leakage of Deterministic Encryption

Mireya Jurado and Geoffrey Smith

November 11, 2019

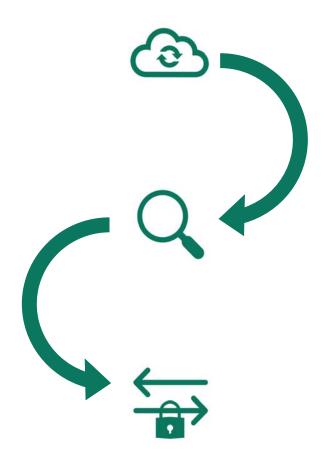
CCSW 2019: Cloud Computing Security Workshop



Challenge

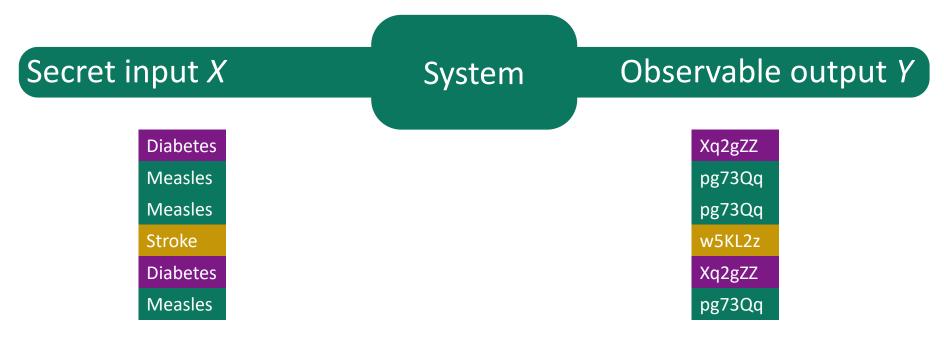
- We want to store sensitive data remotely
- Encrypted databases aim to balance security and functionality





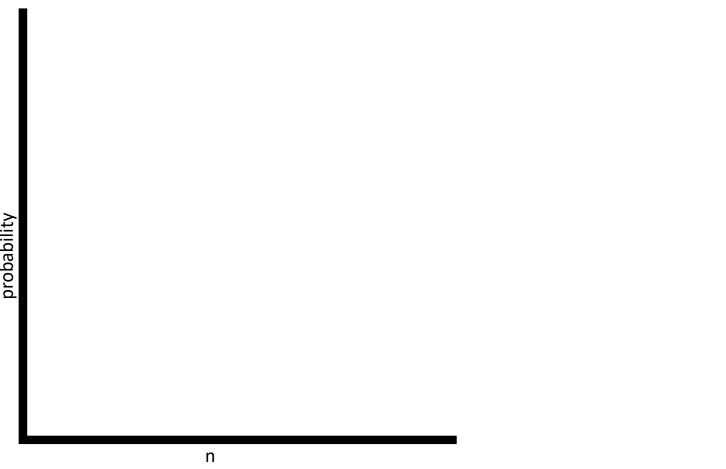
Question: How do you measure information leakage of deterministic encryption?

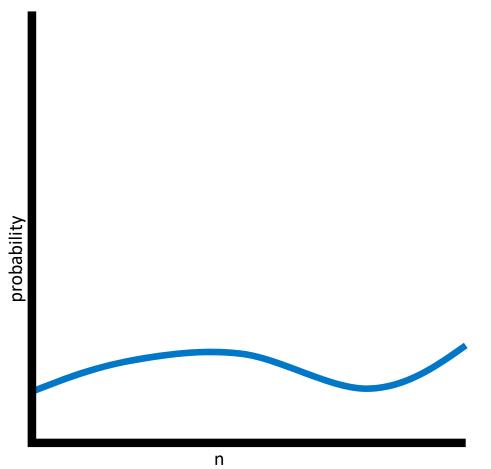
Quantitative Information Flow (QIF)



The database column of diseases, drawn independently according to some prior distribution δ on diseases

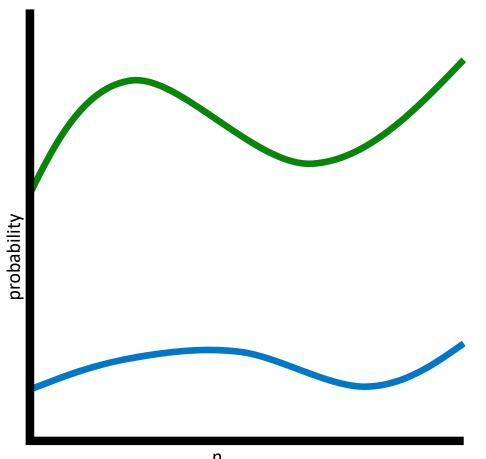
The deterministic encryption of the column, modeled as a random permutation (the "ideal object")





Prior vulnerability:

Adversary's probability of accomplishing goal given only δ

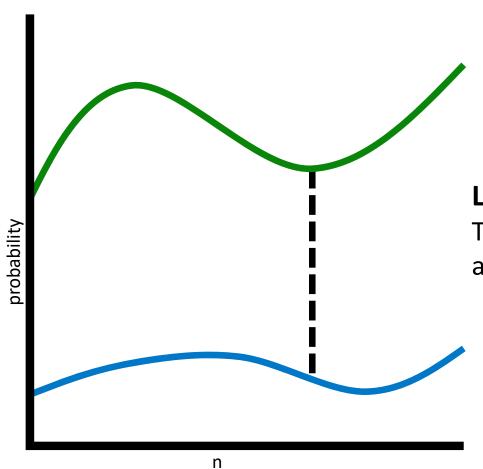


Posterior vulnerability:

Adversary's probability of accomplishing goal given δ and output Y

Prior vulnerability:

Adversary's probability of accomplishing goal given only δ



Posterior vulnerability:

Adversary's probability of accomplishing goal given δ and output Y

Leakage:

The difference between prior and posterior vulnerability

Prior vulnerability:

Adversary's probability of accomplishing goal given only δ

Intuition of Model

Disease	Probability δ
а	1/2
b	1/3
С	¹ / ₆

Xq2gZZ pg73Qq pg73Qq w5KL2z Xq2gZZ pg73Qq

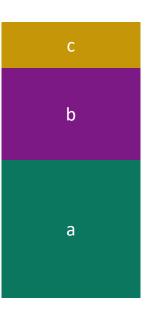
Intuition of Model

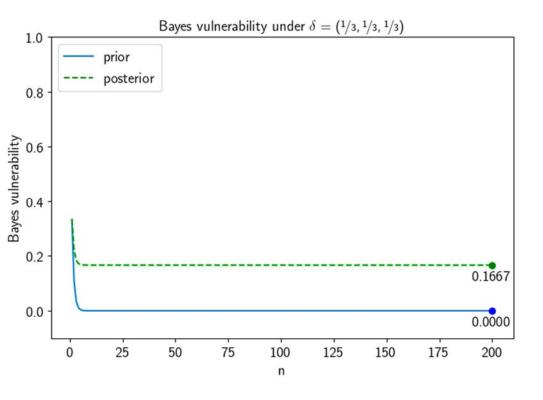
Disease	Probability δ
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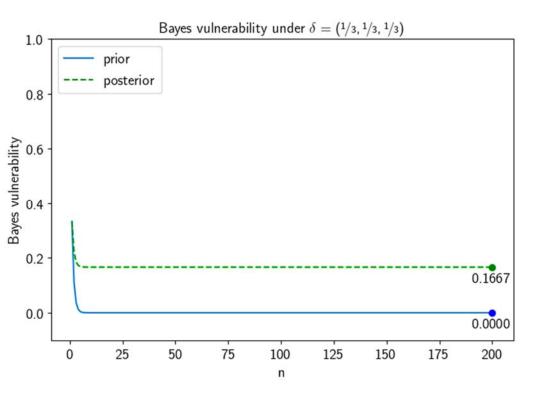
Intuition of Model

Disease	Probability δ
а	1/2
b	1/3
С	1/6





Goal: guess the entire secret in one try

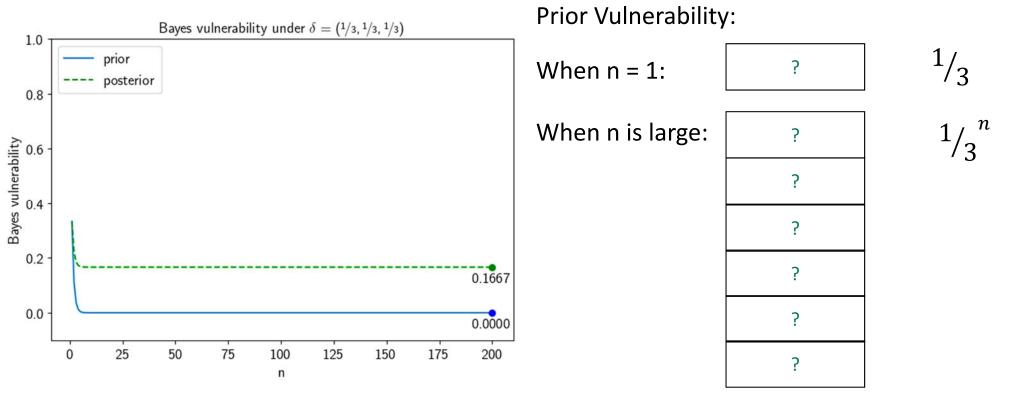


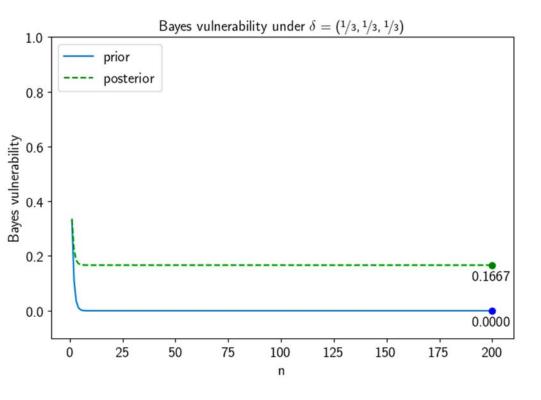
Prior Vulnerability:

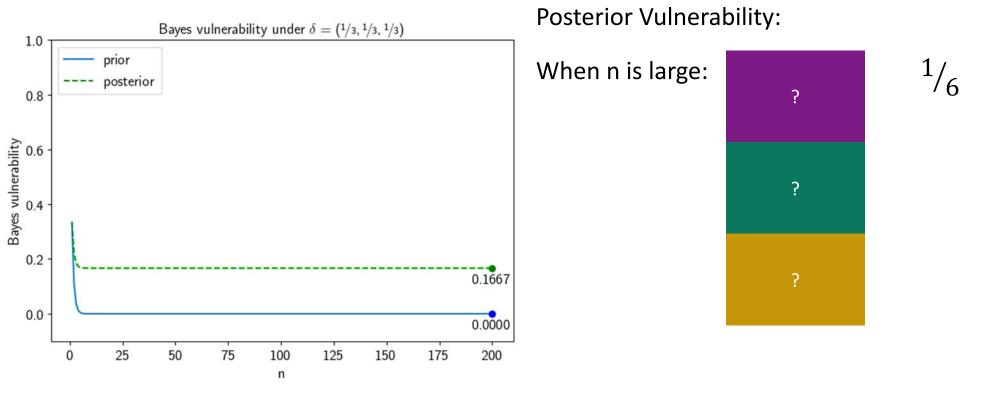
When n = 1:

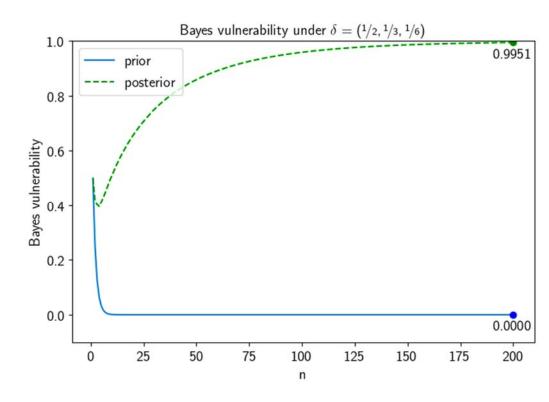
?

1/5





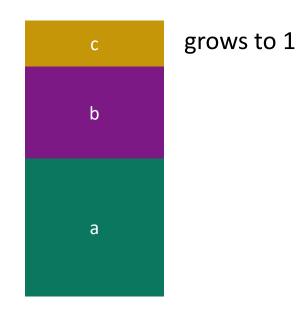


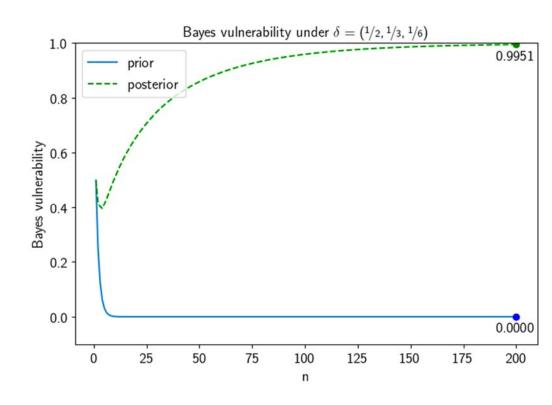


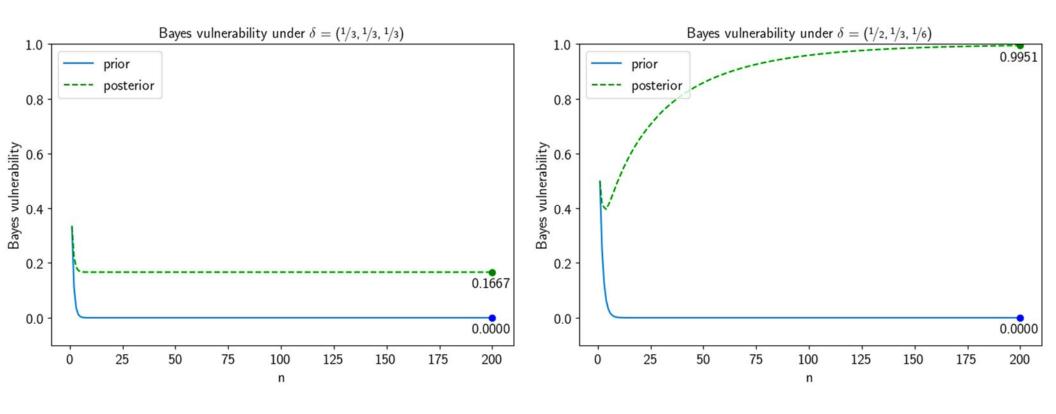
Goal: guess the entire secret in one try

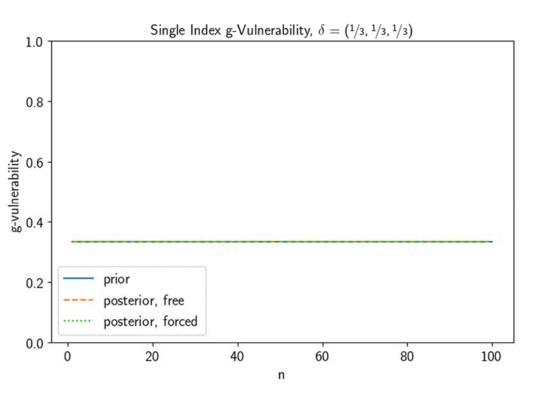
Posterior Vulnerability:

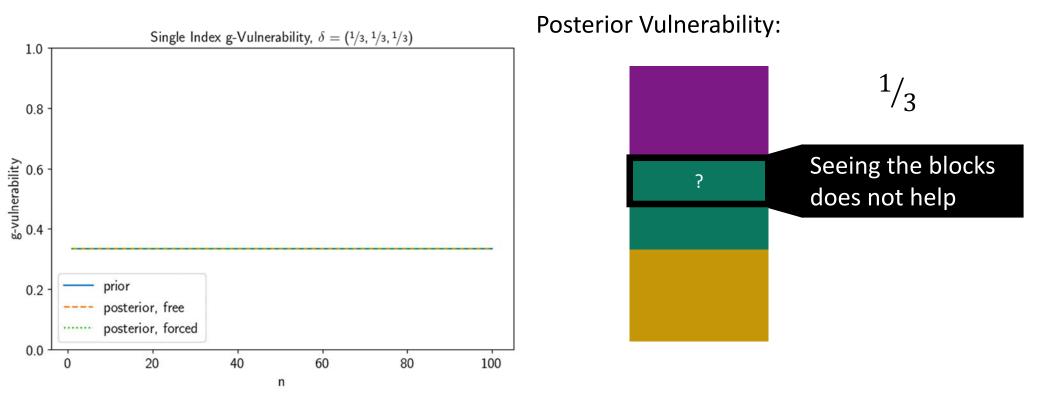
When n is large:

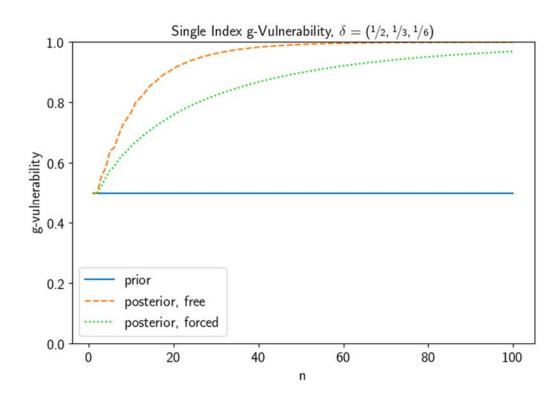






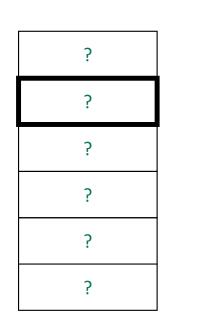




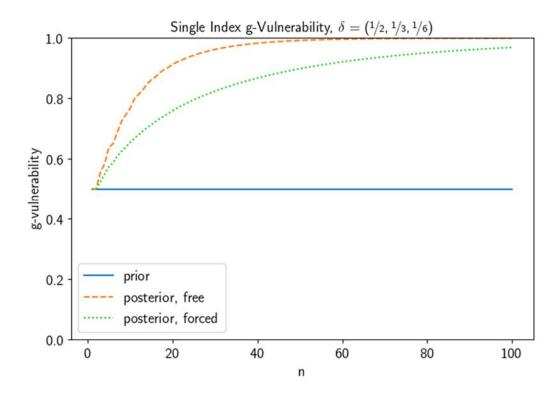


Goals: (1) free to guess any patient's disease and (2) forced to guess a specified patient's disease

Prior Vulnerability:

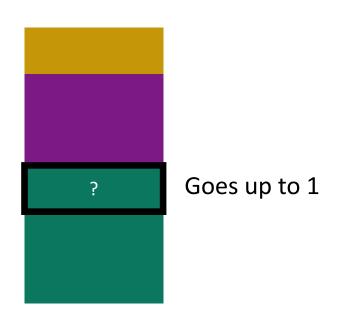


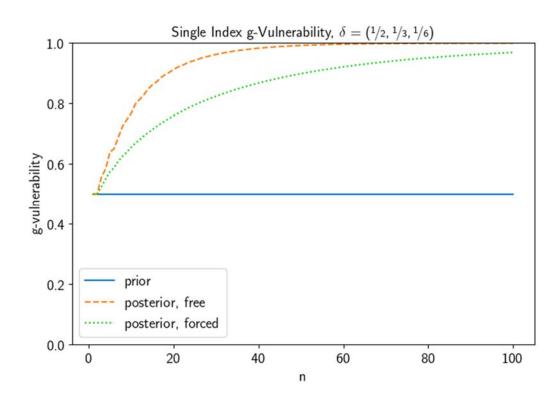
 $^{1}/_{2}$



Goals: (1) free to guess any patient's disease and (2) forced to guess a specified patient's disease

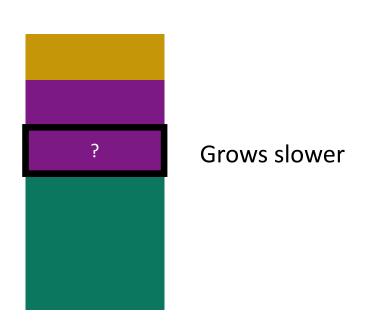
Posterior Vulnerability (free):

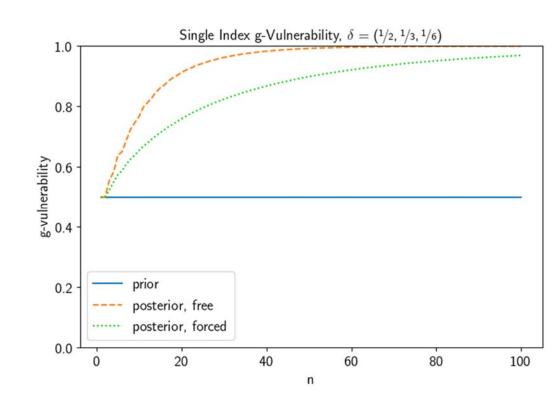




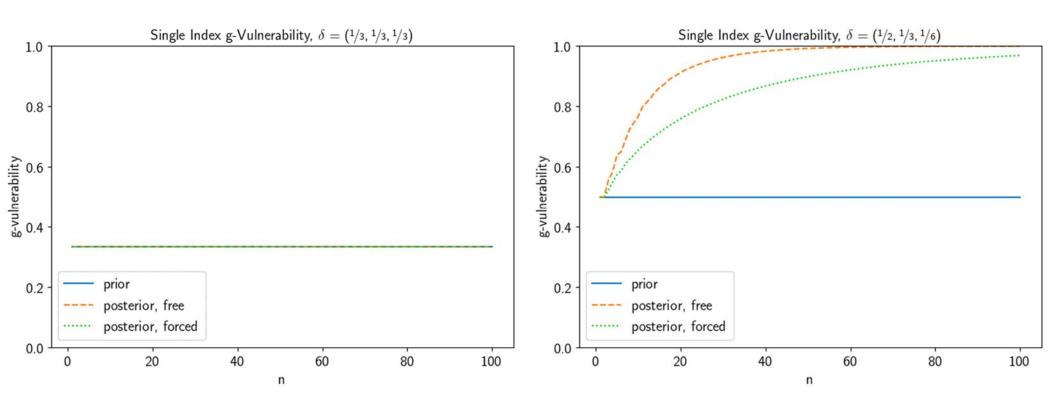
Goals: (1) free to guess any patient's disease and (2) forced to guess a specified patient's disease

Posterior Vulnerability (forced):



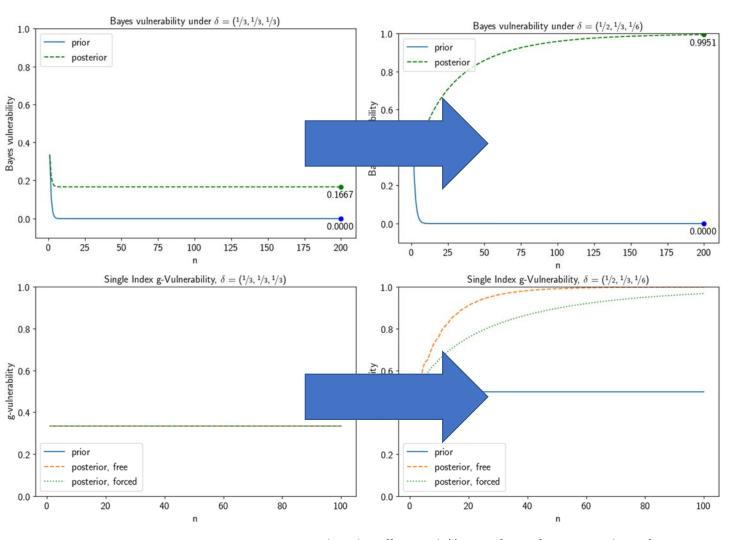


Goals: (1) free to guess any patient's disease and (2) forced to guess a specified patient's disease

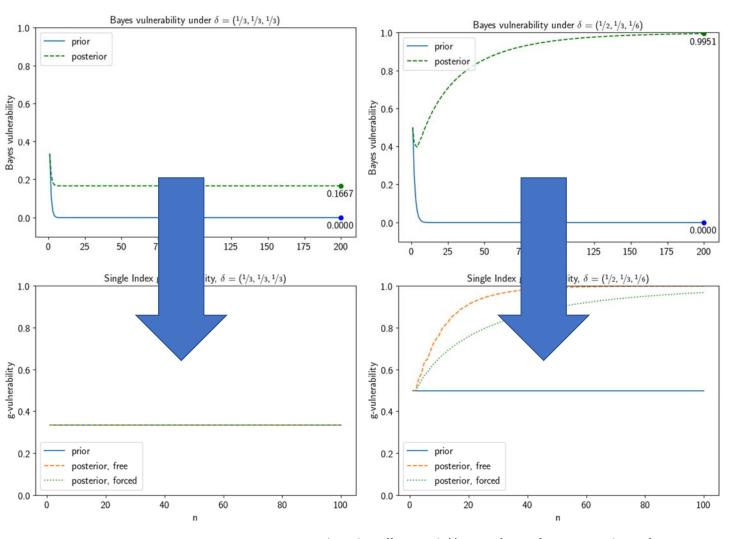


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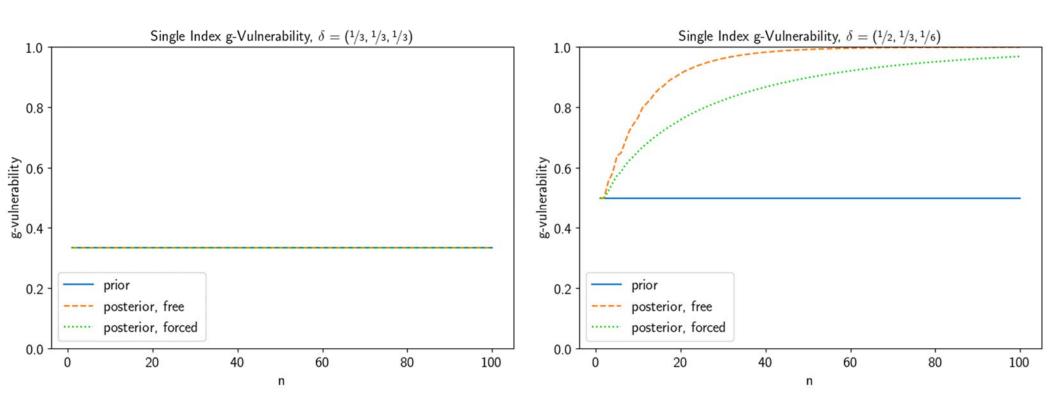
Leakage
Depends on
Prior and
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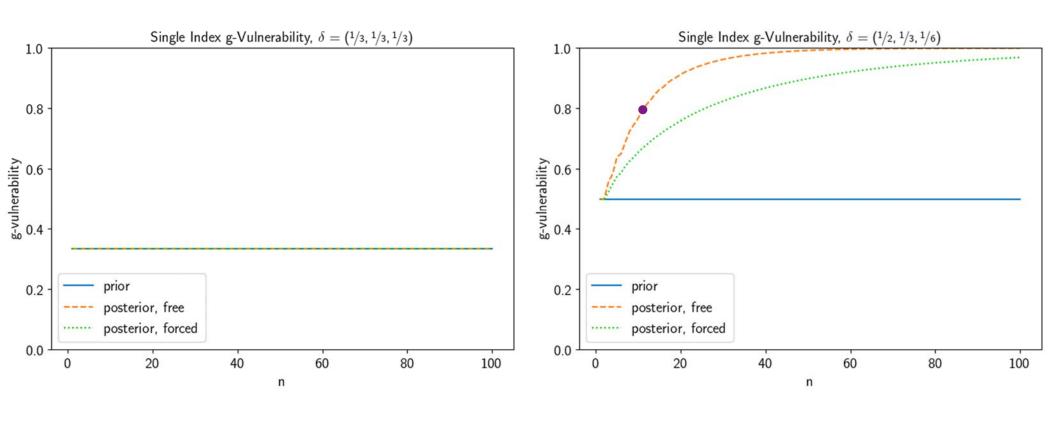
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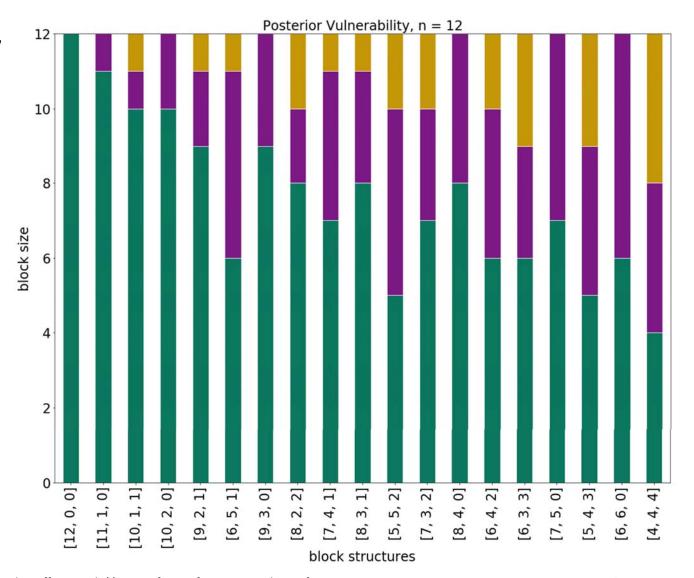
8 of 10

Goals: (1) free to guess any patient's disease and (2) forced to guess a specified patient's disease

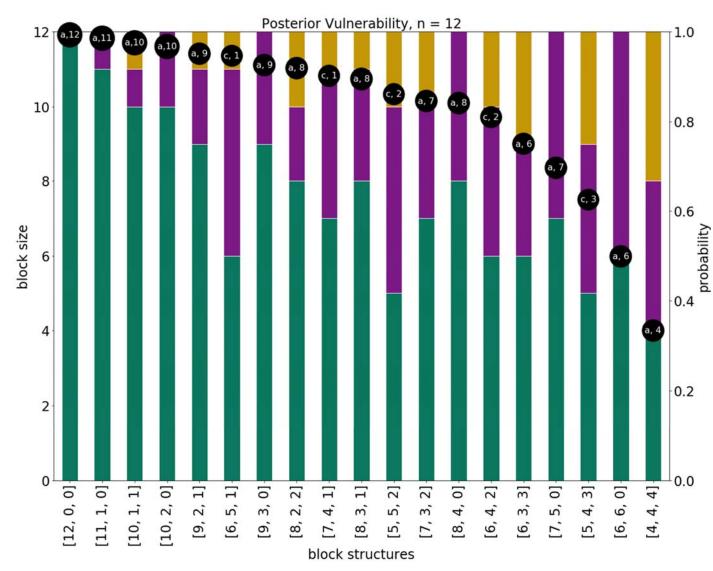


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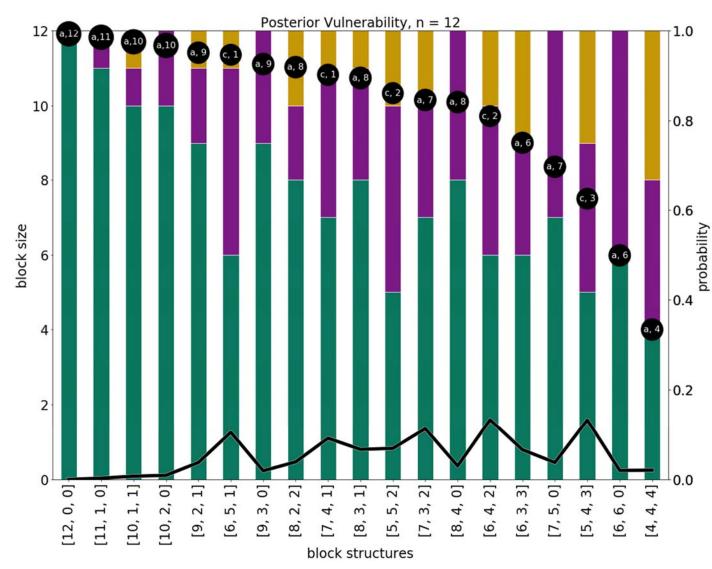
19 possible block structures



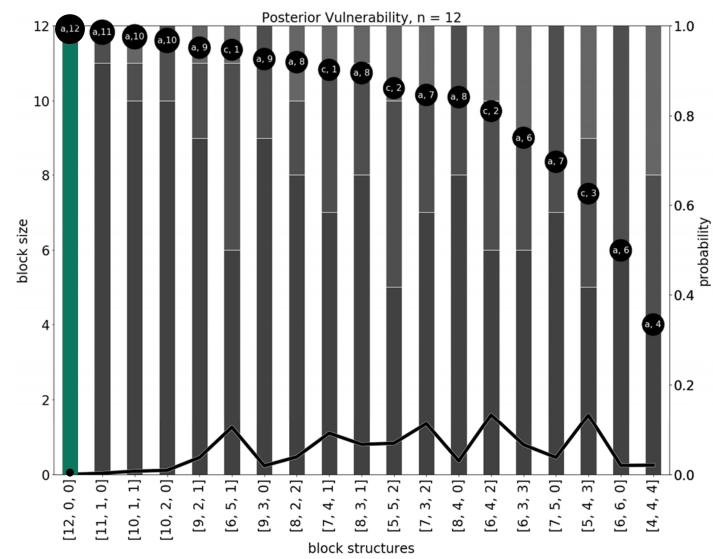
- 19 possible block structures
- Dots: the best guess and her probability of being correct



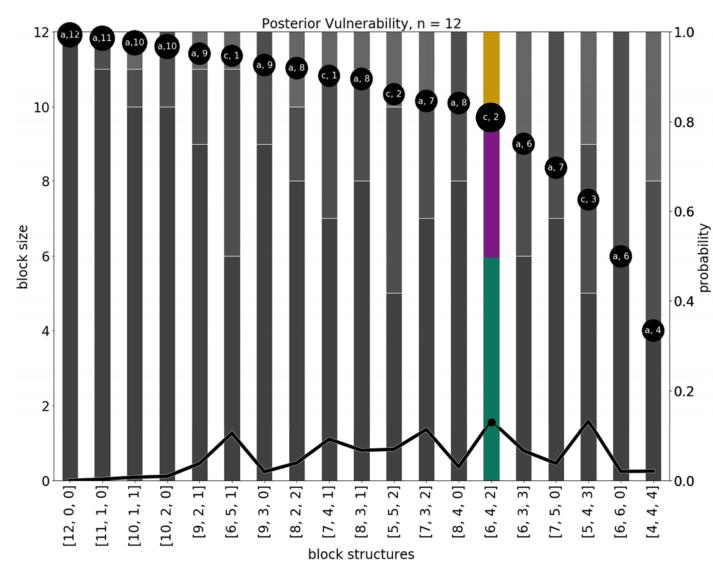
- 19 possible block structures
- Dots: the best guess and her probability of being correct
- Line: the probability the block structure will occur



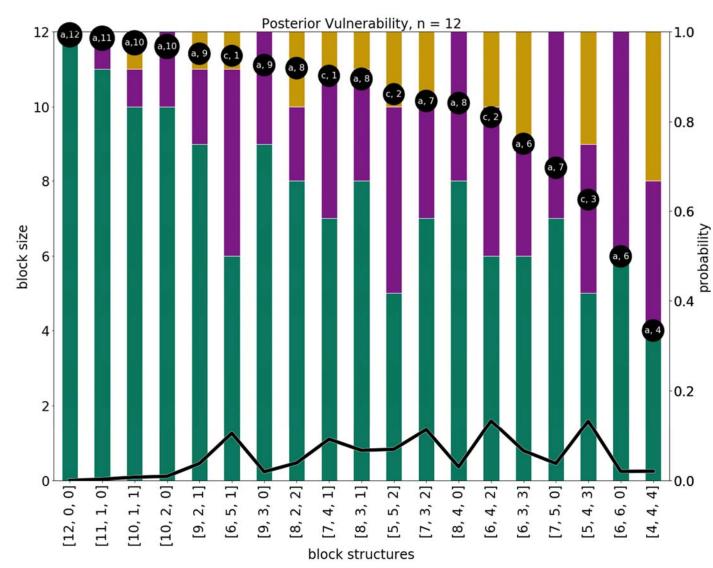
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- 19 possible block structures
- Dots: the best guess and her probability of being correct
- Line: the probability the block structure will occur
- Posterior vulnerability: ~0.813



Contributions



Revealed that there are scenarios where deterministic encryption is not safe even under a uniform prior.



Demonstrated a novel security framework by coupling the provable security approach of modern cryptography with QIF theory.



Illustrated that there is no one 'right' way to quantify leakage. Information flow depends on the operational scenario.